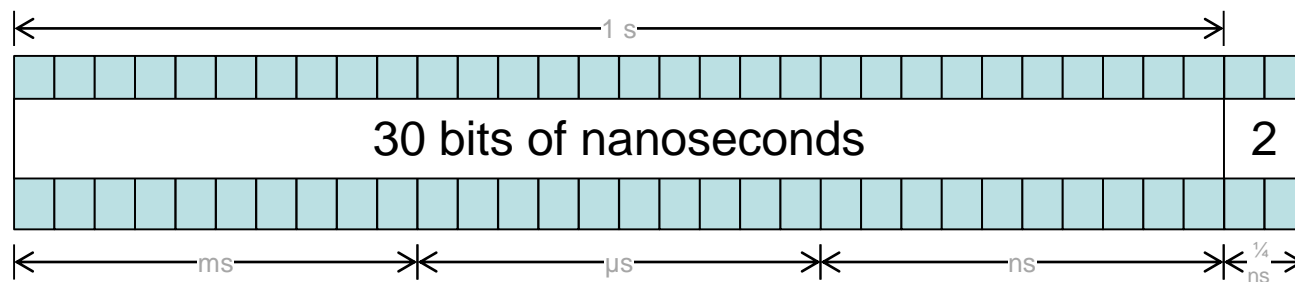




Timestamp Precision

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- The *orderInfo* field can be used as a 32-bit timestamp, down to $\frac{1}{4}$ ns granularity:



- Two main uses of timestamp:
 - Indicating start or end time of flow
 - Indicating presentation time of packets for flows with non-constant data rates

s = second
ms = millisecond
μs = microsecond
ns = nanosecond

- ❑ To reduce bandwidth during idle periods, some protocols will have variable rates
 - Fronthaul may be variable, even if rate to radio unit itself is a constant rate
- ❑ Presentation times allows RoE to handle variable data rates
 - Data may experience jitter in network
 - Egress buffer compensates for network jitter
 - Presentation time is when the data is to exit the RoE node
 - Jitter cleaners ensure data comes out cleanly, and on the right bit period

Jitter vs. Synchronization

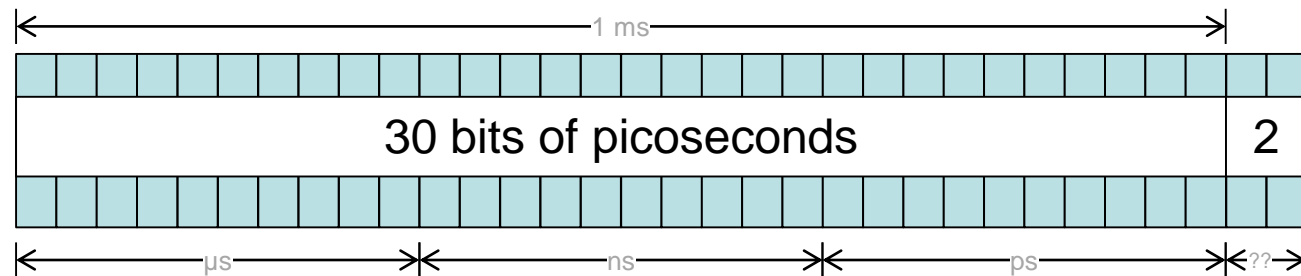
- ❑ Synchronization requirements for LTE are *only* down to $\sim \pm 65$ ns accuracy
 - Each RoE node may be off from TAI by up to 65 ns (or more in some circumstances)
 - Starting and ending a stream may be off by this amount
- ❑ ...but jitter from packet to packet must be much tighter
 - RoE nodes should be able to output data at precise relative times if timestamp is used for a given packet
 - Relative bit time within a flow is important

- ❑ Is the current $\frac{1}{4}$ ns granularity tight enough for today's systems, and does it have headroom for the future?
 - Each bit in 9.8 Gbps CPRI is $\sim 1/10$ ns
 - Each bit in 24 Gbps CPRI is $\sim 1/24$ ns
 - Rates of 100 Gbps or more are likely in the reasonable future

- ❑ How do you specify a presentation time with bit times that may be tiny fractions of nanoseconds if the smallest unit is in $\frac{1}{4}$ nanoseconds?

- ❑ Assume 100 Gbps raw data rate, with extended idle periods suppressed
- ❑ Raw data:
 - ..., 0x3F, 0x4E, <807 bytes of 0's>, 0x39, 0x41
 - One packet ends with 0x4E
 - Next packet starts with 0x99
- ❑ How does RoE say when that packet is supposed to hit?
 - One bit position early, first byte = 0x72
 - One bit position late, first byte = 0x1C
- ❑ Relative timing of bits is important

- Define timestamps in terms of picoseconds, rather than nanoseconds:

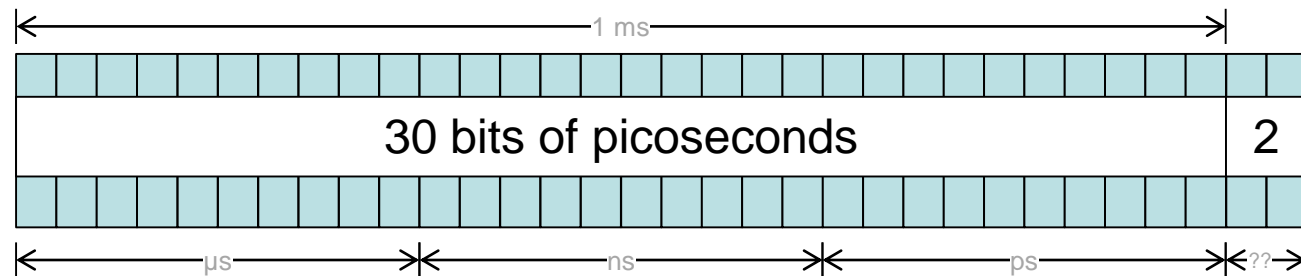


Justification:

- Prior discussions indicated transit time will be less than 1 ms now and in future
- Provides headroom for future speed increases
- No change in field length or usage

ms = millisecond
 μs = microsecond
ns = nanosecond
ps = picosecond

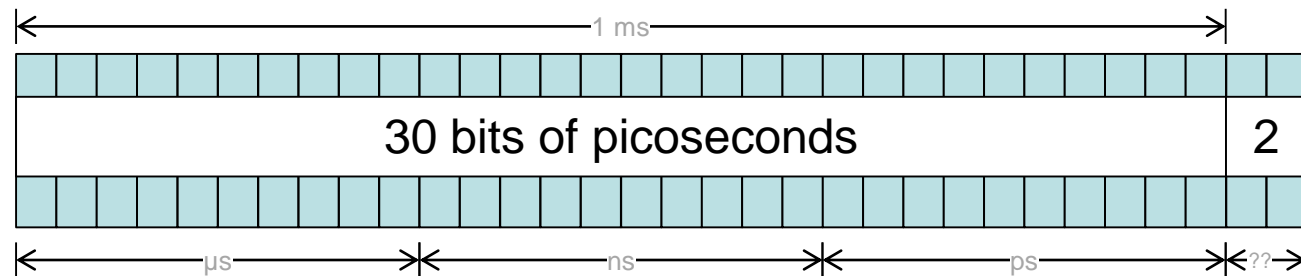
- 30 bits will handle down to picosecond level—what about the last 2 bits?



- Two options (the first is recommended):
 - Implement mini sequence number (0, 1, 2, 3, 0, 1, 2, 3, ...) to detect missed packets
 - Get some benefits of sequence number w/timestamp
 - ... or continue down to $\frac{1}{4}$ ps time

ms = millisecond
μs = microsecond
ns = nanosecond
ps = picosecond

- This presentation proposes re-defining timestamp to handle picosecond timing:



- This handles presentation times up to 1 ms in the future, while offering precision to 1 ps
- Lower 2 bits could implement 2-bit sequence number to detect the occasional missing packet
 - Or could be used to go down to $\frac{1}{4}$ ps timing
- Timestamp purpose/usage unchanged

ms = millisecond
 μ s = microsecond
ns = nanosecond
ps = picosecond